

# Discrete-Event Simulation: A First Course

## Section 3.2: Multi-Stream Lehmer RNGs

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- Typical DES models have many stochastic components
- Want a unique source of randomness for each component
- One (poor) option: multiple RNGs
- Better option: one RNG with multiple “streams” of random numbers

one stream per stochastic component

- We will partition output from our Lehmer RNG into multiple streams

## Example 3.2.1: ssq2 Arrival and Service Processes

- ssq2 has two stochastic components: arrival and service
- Allocate a different generator state variable to each

### GetService with Unique Seed

```
double GetService(void)
{
    double s;
    static long x = 12345;
    PutSeed(x);
    s = Uniform(1.0, 2.0);
    GetSeed(&x);
    return (s);
}
```

- x represents the current state of the service process

## Example 3.2.2: ssq2 Arrival and Service Processes

- Arrival should have its own static variable, initialized differently

### GetArrival with Unique Seed

```
double GetArrival(void)
{
    static double arrival = START;
    static long x = 54321;
    PutSeed(x);
    arrival += Exponential(2.0);
    GetSeed(&x);
    return (arrival);
}
```

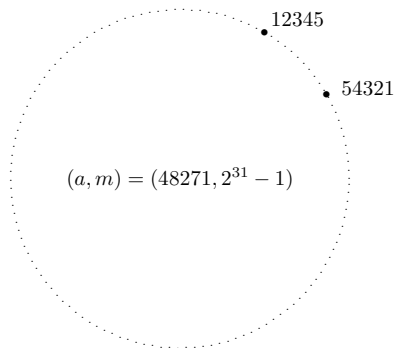
- $x$  represents the current state of arrival process

# The Modified Arrival and Service Processes

- As modified, arrival and service times are drawn from different streams of random numbers
- Nothing magic about the choice of seed for each stream
- The choices may, in fact, be poor ones!
- Provided the streams don't overlap, the processes are *uncoupled*
- Execution time cost is negligible (see Example 3.2.3)

# Stream Considerations

- Potential problem: assignment of initial seeds to facilitate streams
- Each initial state should be chosen to produce *disjoint* streams
- If states are picked at whim, no guarantee of disjoint streams
- Some initial states may only be a few calls to Random apart



# Jump Multipliers

- We will develop a multi-stream version of rng

## Theorem (3.2.1)

Given  $g(x) = ax \bmod m$  and integer  $j (1 < j < m - 1)$

*Jump function:*  $g^j(x) = (a^j \bmod m)x \bmod m$

*Jump multiplier:*  $a^j \bmod m$

If  $g(\cdot)$  generates  $x_0, x_1, x_2, \dots$  then  $g^j(\cdot)$  generates  $x_0, x_j, x_{2j}, \dots$

- Theorem 3.2.1 is the key to creating streams

## Example 3.2.4: An Example Jump Function

- If  $m = 31$  and  $a = 3$  and  $j = 6$ , the jump multiplier is

$$a^j \bmod m = 3^6 \bmod 31 = 16$$

- If  $x_0 = 1$  then  $g(x) = 3x \bmod 31$  generates

1, 3, 9, 27, 19, 26, 16, 17, 20, 29, 25, 13, 8, 24, 10, 30, 28, 22, 4, . . .

- The jump function  $g^6(x) = 16x \bmod 31$  generates

1, 16, 8, 4, 2, . . .

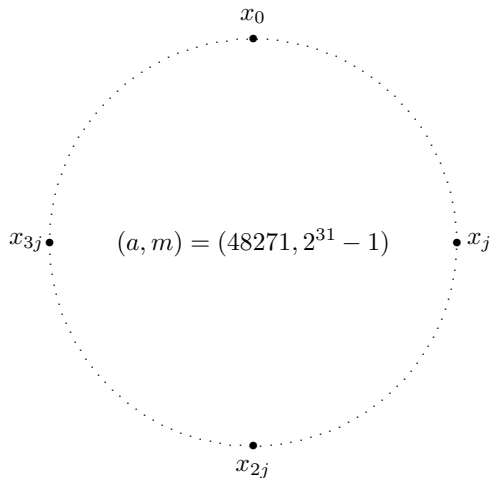
- I.e., the first sequence is  $x_0, x_1, x_2, \dots$ ; the second is  $x_0, x_6, x_{12}, \dots$



# Using the Jump Function

- First, compute the jump multiplier  $a^j \bmod m$  (one time cost)
- Then,  $g^j(\cdot)$  permits jumping from  $x_0$  to  $x_j$  to  $x_{2j}$  to ...
- The user supplies *one* initial seed
- If  $j$  is chosen well,  $g^j(\cdot)$  can “plant” additional initial seeds
- Each planted seed corresponds to a different stream
- Each planted seed is separated by  $j$  calls to Random

# An Example 4-Stream Sequence



## Example 3.2.5: An Appropriate Jump Multiplier

- Consider  $256 = 2^8$  different streams of random numbers
- Partition the RNG output sequence into 256 disjoint subsequences of equal length
- Find the largest  $j < 2^{31}/2^8 = 2^{23}$  such that the jump multiplier is modulus-compatible
- $g^j(x) = (48271^j \bmod m)x \bmod m$  can be implemented via Alg 2.2.1
- Then  $g^j(x)$  can be used to plant the other 255 initial seeds
- Possibility of stream overlap is minimized (though not eliminated!)

# Maximal Modulus-Compatible Jump Multipliers

- *Maximal jump multiplier*:  $a^j \bmod m$  where  $j$  is the largest integer less than  $\lfloor m/s \rfloor$  such that  $a^j \bmod m$  is modulus compatible
- **Example 3.2.6**: multipliers for  $(a, m) = (48271, 2^{31} - 1)$  RNG

# of streams $s$	$\lfloor m/s \rfloor$	jump size $j$	jump multiplier $a^j \bmod m$
1024	2097151	2082675	97070
512	4194303	4170283	44857
256	8388607	8367782	22925
128	16777215	16775552	40509

# Library rngs

- `rngs` is an upward-compatible multi-stream replacement for `rng`
- By default, provides 256 streams, indexed 0 to 255 (0 is the default)
- Only one stream is active at any time
- Six available functions:
  - `Random(void)`
  - `PutSeed(long x)`: superseded by `PlantSeeds`
  - `GetSeed(long *x)`
  - `TestRandom(void)`
  - `SelectStream(int s)`: used to define the active stream
  - `PlantSeeds(long x)`: “plants” one seed per stream
- Henceforth, `rngs` is the library of choice

## Example 3.2.7: ssq2 Revisited

- Use rngs functions for GetArrival, GetService

### GetArrival Method

```
double GetArrival(void) {  
    static double arrival = START;  
    SelectStream(0);  
    arrival += Exponential(2.0);  
    return (arrival);  
}
```

### GetService Method

```
double GetService(void) {  
    SelectStream(2);  
    return (Uniform(1.0, 2.0));  
}
```

- Include "rngs.h" and use PlantSeeds(12345)

# Uncoupling Stochastic Processes

- Per modifications, arrival and service processes are uncoupled
- Consider changing the service process to

$\text{Uniform}(0.0, 1.5) + \text{Uniform}(0.0, 1.5)$

- Without uncoupling, arrival process sequence would change!
- With uncoupling, the service process “sees” *exactly* the same arrival sequence
- Important variance reduction technique

# Single-Server Service Node with Multiple Job Types

- Extend the single-server service node model from Chapter 1
- Consider multiple job types, each with its own arrival and service process
- **Example 3.2.8:** Suppose there are two job types
  - 1 *Exponential*(4.0) interarrivals, *Uniform*(1.0, 3.0) service
  - 2 *Exponential*(6.0) interarrivals, *Uniform*(0.0, 4.0) service

Use `rngs` to allocate a different stream to each stochastic process



# Example 3.2.8: Arrival Process

## Arrival Process

```
double GetArrival(int *j)
/* Index j corresponds to job type */
{
    const double mean[2] = {4.0, 6.0};
    static double arrival[2] = {START, START};
    static int init = 1;
    double temp;
    if (init) {
        SelectStream(0);
        arrival[0] += Exponential(mean[0]);
        SelectStream(1);
        arrival[1] += Exponential(mean[1]);
        init = 0;
    }
    if (arrival[0] <= arrival[1])
        *j = 0;
    else
        *j = 1;
    temp = arrival[*j];
    SelectStream(*j);
    arrival[*j] += Exponential(mean[*j]);
    return (temp);
}
```

## Example 3.2.8: Service Process

### Service Process

```
double GetService(int j)
{
    const double min[2] = {1.0, 0.0};
    const double max[2] = {3.0, 4.0};
    SelectStream(j + 2);
    return (Uniform(min[j], max[j]));
}
```

- Index  $j$  matches service time to appropriate job type
- All four simulated stochastic processes are uncoupled
- Any process could be changed without altering the random sequence of others!

# Consistency Checks

- With appropriate changes to ssq2, steady-state statistics are

$\bar{r}$	$\bar{w}$	$\bar{d}$	$\bar{s}$	$\bar{l}$	$\bar{q}$	$\bar{x}$
2.40	7.92	5.92	2.00	3.30	2.47	0.83

- Obvious consistency checks:  $\bar{w} = \bar{d} + \bar{s}$  and  $\bar{l} = \bar{q} + \bar{x}$
- Other consistency checks:
  - Both job types have avg service time of 2.0  $\implies \bar{s} = 2.00$
  - Net arrival rate should be  $1/4 + 1/6 = 5/12 \implies \bar{r} = 12/5 = 2.40$
  - $\bar{x}$  should be ratio of arrival to service rates

$$\frac{5/12}{1/2} = 5/6 \cong 0.83$$