

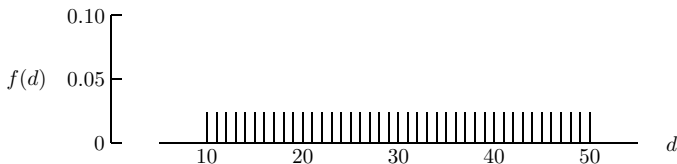
Discrete-Event Simulation: A First Course

Section 6.3: Discrete Random Variable Applications

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Example 6.3.1: The inventory demand model in program `sis2`

- The demand per time interval is an *Equilikely*(10,50) random variate
- $\mu = 30$, $\sigma = \sqrt{140} \cong 11.8$, and the demand pdf is flat



- This model is not very realistic (see Chapter 9)

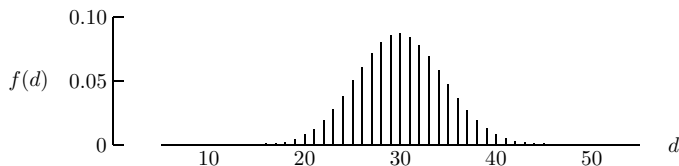
Alternative Inventory Demand Model

- Consider a *Binomial*(100,0.3) model
 - 100 instances per time interval when demand for 1 unit may occur
 - The probability of demand is 0.3 per instance (independently)
 - The function `GetDemand` in `sis2` becomes:

Modified GetDemand Method

```
long GetDemand(void) {
    return (Binomial(100,0.3));
}
```

- $\mu = 30$, $\sigma = \sqrt{21} \cong 4.6$ and the pdf is:



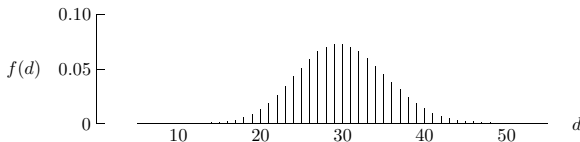
Example 6.3.2: A Poisson(30) Model

- Recall that $Binomial(n, p) \approx Poisson(np)$ for large n
- If $Binomial(100, 0.3)$ is realistic, should also consider $Poisson(30)$
- The function `GetDemand` in program `sis2` would be

Modified GetDemand Method

```
long GetDemand(void) {
    return (Poisson(30.0));
}
```

- $\mu = 30$, $\sigma = \sqrt{30} \cong 5.5$ and the pdf has slightly "heavier" tails



- $Poisson(\lambda)$ is the inventory demand model used in `sis3` with $\lambda = 30$

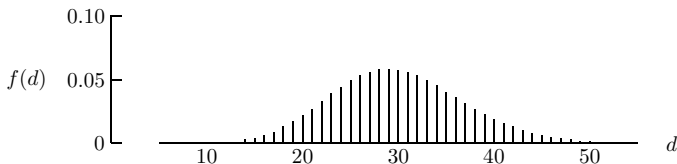
Example 6.3.3: A Pascal(50,0.375) Model

- 50 instances per time interval
- The demand per instance is *Geometric*(p) with $p = 0.375$
- The function `GetDemand` in program `sis2` would be

Modified GetDemand Method

```
long GetDemand(void) {
    return return (Pascal(50,0.375));
}
```

- $\mu = 30$, $\sigma = \sqrt{48} \cong 6.9$ and the pdf has heavier tails than the *Poisson*(30) pdf



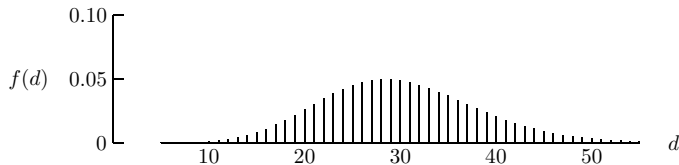
Example 6.3.4

- The number of demand instances per time interval is $Poisson(50)$
- The demand per instance is $Geometric(p)$ with $p = 0.375$

Modified GetDemand Method

```
long GetDemand(void) {
    long instances = Poisson(50.0); /* avoid 0 */
    return (Pascal(instances, 0.375));
}
```

- $\mu = 30$, $\sigma = \sqrt{66} \cong 8.1$ and the pdf has heavier tails



The pdf in Example 6.3.4

- Define random variables

D : the demand *amount*

I : the number of demand *instances* per time interval

$$f(d) = \Pr(D = d) = \sum_{i=0}^{\infty} \Pr(I = i) \Pr(D = d | I = i) \quad d = 0, 1, 2, \dots$$

- To compute $f(d)$, truncate infinite sum: $0 < a \leq i \leq b$

Computing $f(d)$

```
/* use the library rvms */
double sum = 0.0;
for (i = a; i <= b; i++)
    sum += pdfPoisson(50.0,i) * pdfPascal(i,0.375,d);
return sum;
/* sum is f(d) */
```

Program sis4

- Based on sis3 but with a more realistic inventory demand model
- The inter-demand time is an *Exponential*($1/\lambda$) random variate
- Whether or not a demand occurs at demand instances is random with probability p
- To allow for the possibility of more than 1 unit of demand, the demand amount is a *Geometric*(p) random variate
- Expected demand per time interval is

$$\frac{\lambda p}{(1 - p)}$$

Example 6.3.5: The Auto Dealership

- The inventory demand model for `sis4` corresponds to λ customers per week on average
- Each customer will buy
 - 0 autos with probability $1 - p$
 - 1 auto with probability $(1 - p)p$
 - 2 autos with probability $(1 - p)p^2$, etc.
- With $\lambda = 120.0$ and $p = 0.2$, average demand is 30.0

$$30.0 = \frac{\lambda p}{1 - p} = \lambda \sum_{x=0}^{\infty} x(1-p)p^x = \underbrace{\lambda(1-p)p}_{19.2000} + \underbrace{2\lambda(1-p)p^2}_{7.680} + \underbrace{3\lambda(1-p)p^3}_{2.304} + \dots$$

- $\lambda(1 - p) = 96.0$ customers buy 0 autos
- $\lambda(1 - p)p = 19.200$ customers buy 1 auto
- $\lambda(1 - p)p^2 = 3.840$ customers buy 2 autos
- $\lambda(1 - p)p^3 = 0.768$ customers buy 3 autos, etc.

Truncation

- In the previous example, no bound on number of autos purchased
- Can be made more realistic by *truncating* possible values
- Start with random variable X with possible values $\mathcal{X} = \{0, 1, 2, \dots\}$ and cdf $F(x) = \Pr(X \leq x)$
- Want to restrict X to the finite range $0 \leq a \leq x \leq b < \infty$
- If $a > 0$, $\alpha = \Pr(X < a) = \Pr(X \leq a - 1) = F(a - 1)$
- $\beta = \Pr(X > b) = 1 - \Pr(X \leq b) = 1 - F(b)$
- $\Pr(a \leq X \leq b) = \Pr(X \leq b) - \Pr(X < a) = F(b) - F(a - 1)$
Essentially, always true iff $F(b) \cong 1.0$ and $F(a - 1) \cong 0.0$

Specifying truncation points

- If a and b are specified
 - Left-tail, right-tail probabilities α and β obtained using cdf

$$\alpha = \Pr(X < a) = F(a-1) \quad \text{and} \quad \beta = \Pr(X > b) = 1 - F(b)$$

- Transformation is exact
- If α and β are specified
 - Idf can be used to obtain a and b

$$a = F^*(\alpha) \quad \text{and} \quad b = F^*(1 - \beta)$$

- Transformation is not exact because X is discrete

$$\Pr(X < a) \leq \alpha \quad \text{and} \quad \Pr(X > b) < \beta$$

Example 6.3.6

For the $Poisson(50)$ random variable I , determine a, b so that

$$\Pr(a \leq I \leq b) \cong 1.0$$

- Use $\alpha = \beta = 10^{-6}$
- Use `rvms` to compute

Determining a, b

```
a = idfPoisson(50.0, alpha);    /*alpha = 10^-6*/
b = idfPoisson(50.0, 1.0 - beta); /*beta = 10^-6*/
```

- Results: $a = 20$ and $b = 87$
- Consistent with the bounds produced by the conversion:
 $\Pr(I < 20) = \text{cdfPoisson}(50.0, 19) \cong 0.48 \times 10^{-6} < \alpha$
 $\Pr(I > 87) = 1.0 - \text{cdfPoisson}(50.0, 87) \cong 0.75 \times 10^{-6} < \beta$

Effects of Truncation

- Truncating $Poisson(50)$ to the range $\{20, \dots, 87\}$ is insignificant: truncated and un-truncated random variables have (essentially) the same distribution
- Truncation is useful for efficiency:
 - When idf is complex, inversion requires cdf search
 - cdf values are typically stored in an array
 - Small range gives improved space/time efficiency
- Truncation is useful for realism:
 - Prevents arbitrarily large values possible from some variates
- In some applications, truncation is significant
 - Produces a new random variable
 - Must be done correctly

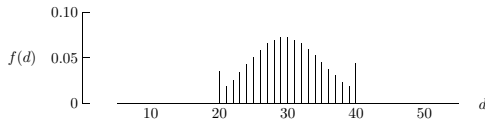
Incorrect Truncation

- Use a $Poisson(30)$ demand model in program `sis2`
- Truncate the demand to the range $20 \leq d \leq 40$

Incorrect Truncation

```
d = Poisson(30.0);
if (d < 20)
    d = 20;
if (d > 40)
    d = 40;
return d;
```

- Original left and right tails grouped together at 20 and 40



- This is *incorrect* for most applications

Truncation by cdf Modification (1)

Example 6.3.8: Truncate $Poisson(30)$ demands to range $20 \leq d \leq 40$

- The $Poisson(30)$ pdf is (before truncation)

$$f(d) = \exp(-30) \frac{30^d}{d!} \quad d = 0, 1, 2, \dots$$

$$\Pr(20 \leq D \leq 40) = F(40) - F(19) = \sum_{d=20}^{40} f(d) \cong 0.945817$$

- Compute a new truncated random variable D_t with pdf $f_t(d)$

$$f_t(d) = \frac{f(d)}{F(40) - F(19)} \quad d = 20, 21, \dots, 40$$

Truncation by cdf Modification (2)

- The corresponding truncated cdf is

$$F_t(d) = \sum_{t=20}^d f_t(t) = \frac{F(d) - F(19)}{F(40) - F(19)} \quad d = 20, 21, \dots, 40$$

- Mean and standard deviation of D_t

$$\mu_t = \sum_{d=20}^{40} df_t(d) \cong 29.841 \quad \text{and} \quad \sigma_t = \sqrt{\sum_{d=20}^{40} (d - \mu_t)^2 f_t(d)} \cong 4.720$$

- Mean and standard deviation of $Poisson(30)$

$$\mu = 30.0 \quad \text{and} \quad \sigma = \sqrt{30} \cong 5.477$$

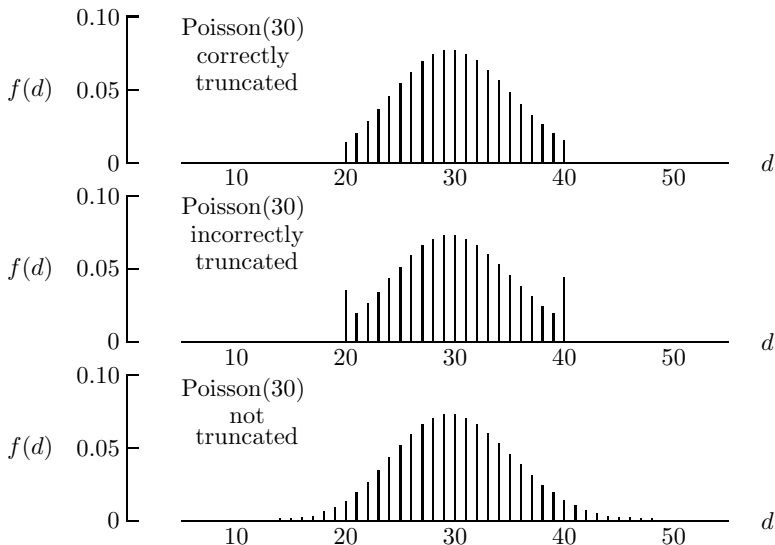
Truncation by cdf Modification (3)

- A random variate truncated to $20 \leq d \leq 40$ can be generated by inversion, using the truncated cdf $F_t(\cdot)$ and Alg.6.2.2

Truncation by cdf Modification

```
u = Random();
d = 30;
if ( $F_t(d) \leq u$ )
    while ( $F_t(d) \leq u$ )
        d++;
else if ( $F_t(20) \leq u$ )
    while ( $F_t(d-1) > u$ )
        d--;
else
    d = 20;
return d;
```

Illustration of pdfs



Truncation By cdf Modification In General

- To truncate (integer-valued, discrete) X to possible values $\mathcal{X}_t = \{a, a + 1, \dots, b\} \subset \mathcal{X}$

$$f_t(x) = \frac{f(x)}{F(b) - F(a - 1)} \quad x \in \mathcal{X}_t$$

$$F_t(x) = \frac{F(x) - F(a - 1)}{F(b) - F(a - 1)} \quad x \in \mathcal{X}_t$$

- Above equations assume $a - 1 \in \mathcal{X}$
- Random values of X_t can be generated using inversion and Alg.6.2.2 with cdf $F_t(\cdot)$

Truncation by Constrained Inversion

- Use the idf of X to generate X_t truncated to $a \leq x \leq b$

Truncation by Constrained Inversion

```

/* assumes a - 1 is a possible value of X */
 $\alpha = F(a-1);$ 
 $\beta = 1.0 - F(b);$ 
 $u = \text{Uniform}(\alpha, 1.0 - \beta);$ 
 $x = F^*(u);$  /*  $F^*(\cdot)$  is the idf of  $X$  */
return x;

```

- The key is that u is *constrained* to a subrange $(\alpha, 1 - \beta) \subset (0, 1)$
- Truncation is automatically enforced prior to inversion

Example 6.3.9

- Generate a $Poisson(30)$ random demand truncated to $20 \leq d \leq 40$

Example 6.3.9

```
 $\alpha$  = cdfPoisson(30.0, 19); /*set-up*/  
 $\beta$  = 1.0 - cdfPoisson(30.0, 40); /*set-up*/  
 $u$  = Uniform( $\alpha$ , 1.0 -  $\beta$ );  
 $d$  = idfPoisson(30.0,  $u$ );  
return  $d$ ;
```

- Uses library `rvms`
- α and β are static variables that are computed once only

Truncation By Acceptance-Rejection

- Truncate $Poisson(30)$ by using *acceptance-rejection*

Truncation By Acceptance-Rejection

```
d = Poisson(30.0);  
while ((d < 20) or ( d > 40))  
    d = Poisson(30.0);  
return d;
```

- Acceptance-rejection is not synchronized or monotone even if the un-truncated generator has these properties
- Truncation by cdf modification or constrained inversion is preferable