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CSCI 454/554 Computer and Network Security

Topic 6. Authentication




Authentication

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- Authentication is the process of reliably verifying certain information.
- Examples
 - User authentication
 - Allow a user to prove his/her identity to another entity (e.g., a system, a device).
 - Message authentication
 - Verify that a message has not been altered without proper authorization.
- A related concept
 - identification

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Identification

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- *Identification* is a process through which one ascertains the identity of another person or entity.
- Authentication and identification are different.
 - Identification requires that the verifier check the information presented against all the entities it knows about,
 - Authentication requires that the information be checked for a single, previously identified, entity.
 - Identification must, by definition, uniquely identify a given entity,
 - Authentication does not necessarily require uniqueness.


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 **Authentication Mechanisms** WILLIAM & MARY

- Password-based authentication
 - Use a secret quantity (the password) that the prover states to prove he/she knows it.
 - Threat: password guessing/dictionary attack


Alice → I'm Alice, the password is *fiddlesticks* → Computer System

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 **Authentication Mechanisms (Cont'd)** WILLIAM & MARY


- Address-based authentication
 - Assume the identity of the source can be inferred based on the network address from which packets arrive.
 - Adopted early in UNIX and VMS
- Berkeley *rtools* (*rsh*, *rlogin*, etc)
 - */etc/hosts.equiv* file
 - List of computers
 - Per user *.rhosts* file
 - List of <computer, account>
- Threat
 - Spoof of network address
 - Not authentication of source addresses

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 **Authentication Mechanisms (Cont'd)** WILLIAM & MARY

- Cryptographic authentication protocols
 - Basic idea:
 - A prover proves some information by performing a cryptographic operation on a quantity that the verifier supplies.
 - Usually reduced to the knowledge of a secret value
 - A symmetric key
 - The private key of a public/private key pair


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Topic 6.1 User Authentication




Authentication and Identity

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- What is **identity**?
 - which characteristics uniquely identifies a person?
 - do we care if identity is unique?
- **Authentication**: verify a user's identity
 - a **supplicant** wishes to authenticate
 - a **verifier** performs the authentication
- What's relationship of identity to **role**, or job function?

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User Authentication Can Be Based On...

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1. What the user **knows**
 - passwords, personal information, a key, a credit card number, etc.
2. What the user **is**
 - Physical characteristics: fingerprints, voiceprint, signature dynamics, iris pattern, DNA, etc.
3. What the user **has** in their possession
 - smart card, (physical) key, smartphone, USB token ...
4. **Where** the user is or can be reached
 - email address, IP address, ...
5. **Who** the user **knows**?

Which of the above is best? Best in what way?

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Crypto-Based Authentication

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- Basic idea: user performs a requested cryptographic operation on a value (a challenge) that the verifier supplies
- Usually based on knowledge of a key (secret key or private key)
- Examples: RSA, zero knowledge proofs, ...
- *We'll look at such protocols in more detail next time*

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Address-Based User Authentication

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- Associates identity with network address or email address
 - used by many web services
- Several early OS functions and tools worked this way
- Benefits? Problems?

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Password Authentication

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Password-Based User Authentication

- User demonstrates knowledge of a secret value to authenticate
 - most common method of user authentication

```

    graph LR
      User[User at computer] -- response --> Server[Server]
      Server -- challenge --> User
  
```

- Threats to password-based authentication?

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Some Issues for Password Systems

- A password should be **easy** to remember but **hard** to guess
 - that's difficult to achieve!
- Some questions
 - what makes a good password?
 - where is the password stored, and in what form?
 - how is knowledge of the password verified?

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Password Storage

- Storing unencrypted passwords in a file is **high risk**
 - compromising the file system compromises all the stored passwords
- Better idea: use the password to compute a one-way function (e.g., a hash, an encryption), and store the **output of the one-way function**
- When user inputs the requested password...
 1. compute its one-way function
 2. compare with the stored value

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Attacks on Passwords

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- Suppose passwords could be up to 9 characters long
- This would produce 10^{18} possible passwords; **320,000 years** to try them all at 10 million a second!
- Unfortunately, not all passwords are equally likely to be used

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Example of a Study

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- In a sample of over 3000 passwords:
 - 500 were easily guessed versions of dictionary words or first name / last name
 - 86% of passwords were easily guessed

Length in characters	1	2	3	4	5	6
Number of passwords	15	72	464	477	706	605 (lower case only)

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Common Password Choices

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- Pet names
- Common names
- Common words
- Dates
- Variations of above (backwards, append a few digits, etc.)

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Dictionary Attacks WILLIAM & MARY

- Attack 1 (online):
 - Create a dictionary of common words and names and their simple transformations
 - Use these to guess the password

Dictionary

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Dictionary Attacks (Cont'd) WILLIAM & MARY

- Attack 2 (offline):
 - Usually F is public and so is the password file
 - In Unix, F is crypt, and the password file is `/etc/passwd`.
 - Compute $F(\text{word})$ for each word in the dictionary
 - A match gives the password

Dictionary

Password file

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Dictionary Attacks (Cont'd) WILLIAM & MARY

- Attack 3 (offline):
 - To speed up search, pre-compute $F(\text{dictionary})$
 - A simple look up gives the password

Dictionary

Pre-computed Dictionary

Password file

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Password Salt

- To make the dictionary attack a bit more difficult
- Salt is a n-bit number between 0 and 2^n
- Derived from, for example, the system clock and the process identifier

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Password Salt (Cont'd)

- Storing the passwords

```

    graph TD
      A[Password + Salt] --> B[F]
      B --> C[F(Password + Salt)]
      C --> D[Username, Salt, F(Password + Salt)]
      D --- E[Password file]
  
```

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Password Salt (Cont'd)

- Verifying the passwords

```

    graph TD
      A[Username, Salt, F(Password + Salt)] -- Fetch Salt according to username --> B[Password + Salt]
      B --> C[F]
      C --> D[F(Password + Salt)]
      A -- Compare --> D
      D --- E[Password file]
  
```

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Does Password Salt Help? WILLIAM & MARY

- Attack 1?
 - Without Salt
 - With Salt

Dictionary

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Does Password Salt Help? WILLIAM & MARY

- Attack 2?
 - Without Salt
 - With Salt

Dictionary

Password file

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Does Password Salt Help? WILLIAM & MARY

- Attack 3?
 - Without Salt
 - With Salt (or change periodically?)

Dictionary

Pre-computed Dictionary

Password file

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Example: Unix Passwords

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- Keyed password hashes are stored, with two-character (16 bit) salt prepended
 - password file is publicly readable
- Users with identical passwords but different salt values will have different hash values

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Password Guidelines For Users

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1. Initial passwords are system-generated, have to be changed by user on first login
2. User must change passwords periodically
3. Passwords vulnerable to a dictionary attack are rejected
4. User should not use same password on multiple sites
5. Be careful to choose the security problems and answers to recover your password
6. Etc.

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Other Password Attacks

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- Technical
 - eavesdropping on traffic that may contain unencrypted passwords (especially keystroke logging)
 - "Trojan horse" password entry programs
 - man-in-the-middle network attack
- "Social"
 - careless password handling or sharing
 - phishing

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The S/Key Protocol

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Using "Disposable" Passwords

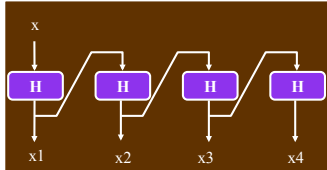
- Simple **idea**: generate a long list of passwords, use each **only one time**
 - attacker gains little/no advantage by eavesdropping on password protocol, or cracking one password
- Disadvantages
 - storage overhead
 - users would have to memorize lots of passwords!
- Alternative: the **S/Key protocol**
 - based on use of **one-way** (e.g. hash) **function**

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S/Key Password Generation

1. Alice selects a password **x**
2. Alice specifies **n**, the number of passwords to generate
3. Alice's computer then generates a sequence of passwords
 - $x_1 = H(x)$
 - $x_2 = H(x_1)$
 - ...
 - $x_n = H(x_{n-1})$



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Generation... (cont'd) WILLIAM & MARY

- 4. Alice communicates (securely) to a server the last value in the sequence: x_n
 - **Key feature:** no one knowing x_i can easily find an x_{i-1} such that $H(x_{i-1}) = x_i$
 - only Alice possesses that information

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Authentication Using S/Key WILLIAM & MARY

- **Assuming** server is in possession of $x_i...$

```

graph LR
    S[Server] -- i --> A[Alice]
    A -- x_{i-1} --> S
    S -- verifies H(x_{i-1}) = x_i --> S
    
```


Is dictionary attack still possible?

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Limitations WILLIAM & MARY


- Value of n limits number of passwords
 - need to **periodically regenerate** a new chain of passwords
- Does not authenticate server! Example attack:
 1. real server sends i to fake server, which is masquerading as Alice
 2. fake server sends i to Alice, who responds with x_{i-1}
 3. fake server then presents x_{i-1} to real server

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Biometrics




Biometrics

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- Relies upon physical characteristics of people to authenticate them
- Desired qualities
 1. uniquely identifying
 2. very difficult to forge / mimic
 3. highly accurate, does not vary
 4. easy to scan or collect
 5. fast to measure / compare
 6. inexpensive to implement
- Which of these are concerns for passwords?

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Assessment

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- Convenient for users (e.g., you always have your fingerprints, never have to remember them), but...
 - potentially troubling sacrifice of private information
 - new wounds on your fingers
 - no technique yet has all the desired properties

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 Example Biometric Technologies WILLIAM & MARY

- Signature / penmanship / typing style
- Fingerprints
- Palm geometry
- Retina scan
- Iris scan
- Face recognition
- Voice recognition

 The Scorecard From One Study WILLIAM & MARY

Biometrics	Univer- sality	Unique- ness	Perma- nence	Collect- ability	Perfor- mance	Accept- ability	Circum- vention
Face	H	L	M	H	L	H	L
Fingerprint	M	H	H	M	H	M	H
Hand Geometry	M	M	M	H	M	M	M
Keystroke Dynamics	L	L	L	M	L	M	M
Hand vein	M	M	M	M	M	M	H
Iris	H	H	H	M	H	L	H
Retina	H	H	M	L	H	L	H
Signature	L	L	L	H	L	H	L
Voice	M	L	L	M	L	H	H
Facial Thermogram	H	H	L	H	M	H	H
DNA	H	H	H	L	H	L	L

H=High, M=Medium, L=Low

 Multifactor Authentication WILLIAM & MARY

- If one characteristic is pretty good, two or more characteristics should be better?
- Suppose true positive rate was **AND** of the two, and false positive rate was **OR** of the two...
 - $TP = TP1 * TP2$
 - $FP = 1 - (1-FP1)*(1-FP2)$
- Alternative: combine a biometric technique with passwords

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Authentication Hardware (Tokens)

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Tokens

- A token is a physical device that can be interfaced to the computer, and carries identifying information
- Types
 - **passive** tokens just store information
 - **active** tokens have processors and can perform cryptographic operations
- Examples
 - cards with magnetic strips
 - smart cards
 - USB storage devices
 - RFID tags


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Design Issues for Tokens


- Cost
- Size
- Capabilities
- Robustness
- Resistance to tampering
- Usefulness if stolen / lost

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 **An Example: Time Synchronized Tokens** WILLIAM & MARY


- The token contains:
 - internal clock
 - display
 - a secret key
- Token computes a one-way function of **current time+key**, and displays that
 - this value changes about once per minute
- User reads this value and types it in to authenticate to the server
 - requires that server and token time stays **synchronized**

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 **One-time Password on Smartphone** WILLIAM & MARY

- Integrate physical tokens into smartphone
- Requirements:
 - Security
 - Malicious mobile OS cannot compromise the keying material in the one-time password (OTP) generator
 - It cannot read the OTP
 - Reliability
 - OTP works even if mobile OS crashes
 - Trusted inputs (e.g., clock time) for the OTP generator
 - Trusted display

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 **TrustZone-based Solution** WILLIAM & MARY

- ARM TrustZone Technology
 - Two isolated execution environments
 - Mobile OS cannot access the disk, memory, CPU states of the OTP generator.
 - A secure clock for OTP generator
 - A self-contained display and touchscreen.

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 **Another Example: Alladin eToken** 

API / standards	PKCS#11 v2.01, CAPI (Microsoft Crypto API), Siemens/Infineon APDU commands, PC/SC, X.509 v3 certificate storage, SSL v3, IPsec/IKE
Security Algorithms	RSA 1024-bit / 2048-bit*, DES, 3DES, SHA1
Power source	Battery, 5 year lifetime
LCD	6 characters
Data retention	10 years



 **Summary** 

1. Passwords are by far the most widely used form of authentication, despite numerous problems
2. Biometrics hold promise but are expensive, inconvenient, and compromise privacy
3. Two factor authentication is commonly used for higher security
4. One-time passwords (S/Key) are attractive, especially if combined with hardware
